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THE QR-ALGORITHM

by

Masako Ogura

September 1, 1971



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THE QR-ALGORITHM

by

Masako Ogura

Center for Advanced Computation
University of Illinois at Urbana-Champaign
Urbana, Illinois 61801

September 1, 1971

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ABSTRACT

The implementation of QR-algorithm on ILLIAC IV is described. An ASK subroutine for computing all eigenvalues of a real Hessenberg matrix of order less than or equal to 64 by this algorithm is attached. The QR-transformation consists of the decomposition of the matrix A_k into the product of a unitary matrix Q_k and an upper triangular matrix R_k , and forming A_{k+1} by post-multiplying R_k by Q_k , where $A_1 = A$ is the original matrix. All eigenvalues are either isolated on the diagonal or are eigenvalues of a 2×2 diagonal submatrix as $k \rightarrow \infty$.

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1. INTRODUCTION

An ASK program for finding all eigenvalues of a real Hessenberg matrix of order less than or equal to 64 is written and tested on the B5500 simulator. The QR-algorithm for real Hessenberg matrices described herein is that of Martin et al [2]. The ILLIAC IV computer time required for performing one iteration (computation of A_{k+2} from A_k , refer to 4.1) on a 64 x 64 matrix is approximately 10 millisecond.

The necessary information for using this program is given in Section 2. The test result of this program on a 4 x 4 real Hessenberg matrix is given in Section 3. In Section 4, the outline of the QR-algorithm is given and Section 5 is devoted to the actual programming technique to implement this algorithm on the ILLIAC IV computer. The flow chart and ASK program are attached as Appendices 1 and 2 respectively.

2. USAGE

This program assumes that the given real Hessenberg matrix is stored in the core memory in the straight storage scheme so that each row is stored across the PE's, starting with PEO. The real and imaginary parts of the eigenvalues found are to be stored in the two rows in the PE memory specified by the user. The original matrix is destroyed and replaced by the matrix which results from the QR transformations. The content of ACAR2 and ACAR3 are destroyed since the ACAR3 is used for linkage between the subroutine and main program and ACAR2 is used for passing the address of parameters to the subroutine.

2.1 Calling sequence

Calling sequence for this subroutine is:

CALL HQR (N, A, WR, WI, IT)

where A designates the first row of the matrix and is declared in the main program as

```
A: DATA      a00, a01 , . . . . . ,      a0,N-1, (0.0)M,
               a10, a11 , . . . . . ,      a1,N-1, (0.0)M,
               . . . . .
               . . . . .
               aN-1,0, aN-1,1, . . . . , aN-1,N-1, (0.0)M;
```

where $M = 64 - N$.

N is the size of the matrix declared as

```
N: EQU 64;
```

or

```
DEFINE N = 64 ##;
```

or given as an integer, i.e., 64.

WR and WI are the rows in the PE memory where resulting real and imaginary parts of the eigenvalues respectively are to be stored. They are declared as

```
WR: BLK 1; and WI: BLK 1;
```

IT is the row vector in memory to which the number of iterations required for finding each eigenvalue is to be placed. This is declared as

```
IT: BLK 1;
```


If zero is placed in place of IT, it is to be considered that a user does not want to know the number of iterations required. The CALL macro should be defined in the user's program as:

```

DEFINE CALL &NAME (&PARAMETERS) =
    &IF &SIGN (&MFIELD(&NAME))
        &THEN EXTERNAL &NAME; &FI
    &IF &EMPTY (&PARAMETERS) &THEN &ELSE
        BEGIN BLOCK
            BEGIN USE (63)
            LIST: DATA &PARAMETERS
            END;
            CLC(2);
            SLIT(2) LIST;
        END; &FI
        CLC(3);
        SLIT(3) &NAME;
        EXCHL(3) &ICR ##;

```

2.2 Core storage used

This routine uses 500 words of PE memory for storing instructions. One row is used for storing PE numbers and three additional rows are used for temporary storage. ADBO ~ 32 are also used.

2.3 Constant

EPS (ϵ), the constant which is used to test the convergence (4.1), is taken as 10^{-10} in this program. If a user wants to change the value of this constant, he may insert EPS: DATA (desirable value); in place of

```
EPS: DATA @ - 10; .
```

3. EXAMPLE

A test of this program was made on the B5500 simulator for the matrix:

$$\begin{bmatrix} 5.0 & -2.0 & -5.0 & -1.0 \\ 1.0 & 0.0 & -3.0 & 2.0 \\ 0.0 & 2.0 & 2.0 & -3.0 \\ 0.0 & 0.0 & 1.0 & -2.0 \end{bmatrix}$$

with $\epsilon = 10^{-8}$. The comparison of the eigenvalues obtained by this program to the exact values is given in Table 1. The selection of this small matrix and a relatively large ϵ was made because of the speed of the SSK simulator on the B5500; the execution speed ratio of the simulator to the ILLIAC IV is approximately $1:10^6$.

Table 1

Eignevalues obtained on B5500 simulator	Exact eigenvalues
3.999999997867	4.0
1.000000001066 + 1.99999999928i	1.0 + 2.0i
1.000000001066 - 1.99999999928i	1.0 - 2.0i
-1.000000000000	-1.0

4. QR-ALGORITHM

4.1 Brief outline of the QR-algorithm

The QR-transformation consists of the decomposition of the matrix A_k into the product of a unitary matrix Q_k and an upper triangular matrix R_k , and forming A_{k+1} by post-multiplying R_k by Q_k . Thus

$$A_{k+1} = R_k Q_k \quad \text{where} \quad A_k = Q_k^* R_k, \quad (1)$$

therefore

$$A_{k+1} = Q_k^* A_k Q_k$$

where $A_1 = A$ is the original matrix. It can be shown in general that A_k tends to a form in which $a_{i+1, i}^{(k)} a_{i+2, i+1}^{(k)} = 0$ for $i = 0, 1, \dots, N-3$ as k increases. All eigenvalues are therefore either isolated on the diagonal or they are eigenvalues of a 2×2 diagonal submatrix. The amount of calculations involved in a QR step is greatly reduced if the matrix A is in the Hessenberg (or almost triangular) form. Since there are several stable methods available to reduce a general matrix to this form (ASK program HSBG is written for this purpose), the QR-algorithm is used after such reduction.

In order to achieve rapid convergence, it is essential that the origin shifts be applied and that each shift be close to an eigenvalue of the matrix. The QR-algorithm with shift of an origin s_k is expressed as:

$$A_{k+1} = R_k Q_k + s_k I \quad \text{where} \quad A_k - s_k I = Q_k R_k \quad (2)$$

or in other words

$$A_{k+1} = Q_k^* A_k Q_k.$$

However, even when A_1 is real, some of the eigenvalues may be complex. If the transformation (2) is carried out with a complex value of s_k , A_{k+1} is in general a complex matrix. This deficiency can be overcome by performing two steps of (2) with shifts of s_k and s_{k+1} respectively. Since s_k and s_{k+1} are both real or complex conjugate in this transformation, A_{k+2} should be always real. This transformation is described as

$$A_{k+2} = Q_{k+1}^* Q_k^* A_k Q_k Q_{k+1}$$

$$\text{and } (Q_k Q_{k+1}) (R_{k+1} R_k) = (A_k - s_k I) (A_k - s_{k+1} I) \quad (3)$$

One method of calculating A_{k+2} by (3) would be to form the real matrix $\Gamma = (A_k - s_k I) (A_k - s_{k+1} I)$, computing its unitary-triangular decomposition to obtain $Q_k Q_{k+1}$ and transform A_k by means of this, thus giving A_{k+2} . This process requires a prohibitive amount of work, but it is shown [1] that when the matrix is in the Hessenberg form, it is unnecessary to compute more than the first column of Γ , and that this immediately gives the transformation to be applied to A_k .

4.2 Practical computation

If (3) is rewritten as

$$A_{k+2} = W^* A_k W \quad \text{and} \quad W^* \Gamma = \Delta$$

where $W = Q_k Q_{k+1}$ and Δ is the triangular matrix $R_{k+1} R_k$, W^* is a unitary matrix which reduced Γ to the triangular Δ , and W is composed of N unitary factors of the form $M_i = \begin{bmatrix} I_i - 1 & 0 \\ 0 & B_i \end{bmatrix}$ so that $W = M_1 M_2 \dots M_k$. From the form of each M_i we see that the first column of W is equal to the first column of M_1 , and this is any unitary matrix, the transpose of which eliminates the elements of the first column of Γ below diagonal.

Since we wish to transform A_k to A_{k+2} by W , we first operate on A_k with M_1 . This will change the first three rows and columns of A_k since the first column of Γ contains only three non-zero elements. It follows;

5. PROGRAM DESCRIPTION

5.1 Search for negligible subdiagonal elements

We assume that the size of the matrix under consideration is $(n + 1) \times (n + 1)$ where n takes integer values between 1 and $N - 1$. If the last negligible subdiagonal element is in position $(\ell, \ell - 1)$, it is required only to work on the submatrix in the rows and columns ℓ to n . If none of the subdiagonal elements are negligible, ℓ is taken to be 0. The following criterion is used

$$|a_{\ell, \ell - 1}| \leq \epsilon (|a_{\ell - 1, \ell - 1}| + |a_{\ell, \ell}|).$$

This criterion examines whether $a_{\ell, \ell - 1}$ is negligible compared to the local diagonal elements.

On each PE ℓ , the following computations are simultaneously performed:

$$f(\ell) \equiv |a_{\ell, \ell - 1}| - \epsilon (|a_{\ell - 1, \ell - 1}| + |a_{\ell, \ell}|)$$

for $\ell = 1, 2, \dots, n$.

If $f(\ell)$ is negative, 1 is placed in ℓ th bit of the ACAR. Then searching is made for the lowest bit of the ACAR which contains 1. For example, in the following case:

	0	1	2	3	...	19	20	63
ACAR	0	0	1	0	...	0	1	0	0	0	...	0	0	0	0	0	0

ℓ is set to be 20.

Then test is made if $\ell = n$ or $\ell = n - 1$. If $\ell = n$, one eigenvalue is found in the place (n, n) and the matrix is deflated by 1, and n is decreased by one. In the case that $\ell = n - 1$, two eigenvalues are found as the eigenvalues of the bottom-right hand corner 2×2 submatrix. Then two columns and rows are deflated and n is decreased by 2.

5.2 Shifts of origin

The shifts of origin at each stage are taken to be the two roots, s_k and s_{k+1} , of the 2×2 matrix in the bottom right-hand corner of the current A_k . This gives

$$s_k + s_{k+1} = a_{n-1, n-1} + a_{n, n} \quad (4)$$

and $s_k s_{k+1} = a_{n-1, n-1} a_{n, n} - a_{n-1, n} a_{n, n-1}$.

In some rare cases, the process fails to converge with these shifts of origin. An example of such failure is provided by matrices of the type:

$$\begin{bmatrix} 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 0 & 0 & 0 \\ & 1 & 0 & 0 & 0 \\ & & 1 & 0 & 0 \\ & & & 1 & 0 \end{bmatrix}.$$

Here the shifts of origin, given in (4), are both zero, and since the matrix is orthogonal, it is invariant with respect to the QR transformation without shifts. However, if one iteration is performed with any shifts of origin which are loosely related to the norm of the matrix, the convergence is very rapid. Therefore in the case where ten iterations do not produce an eigenvalue, the usual shifts s_k and s_{k+1} are replaced by shifts defined by

$$\begin{aligned} s_k + s_{k+1} &= 1.5 (|a_{n, n-1}| + |a_{n-1, n-2}|) \\ s_k s_{k+1} &= (|a_{n, n-1}| + |a_{n-1, n-2}|)^2 \end{aligned} \quad (5)$$

This strategy is used again after 20 unsuccessful iterations. If 30 unsuccessful iterations are needed then a failure indication is given.

In this program, ITS is the name of an ADB where the iteration count is stored. When $ITS \neq 10, 20$, we form $S = s_k + s_{k+1}$ and $Y = s_k s_{k+1}$ on a PE n according to (4) and store in ADB's. When $ITS = 10$ or 20 , scheme (5) is used for computing S and Y . If $ITS = 30$, it is assumed that this algorithm fails to produce eigenvalues. As a result, only eigenvalues computed prior to this point are given in WR and WI.

5.3 Search for two consecutive small subdiagonal elements

After determining ℓ , (5.1), the submatrix in the rows ℓ to n are examined to see if any two consecutive subdiagonal elements are small enough to work with an even smaller submatrix. To test if we are to start at the row m , we compute the elements p_m , q_m and r_m such that

$$\begin{aligned} p_m &= a_{mm}^2 - a_{mm} (s_k + s_{k+1}) + s_k s_{k+1} + a_{m,m+1} a_{m+1,m} \\ q_m &= a_{m+1,m} (a_{mm} + a_{m+1,m+1} - s_k - s_{k+1}) \\ r_m &= a_{m+2,m+1} a_{m+1,m} \end{aligned} \quad (6)$$

The criterion applied is

$$\begin{aligned} |a_{m,m-1}| (|q_m| + |r_m|) \\ \leq \epsilon |p_m| (|a_{m+1,m+1}| + |a_{m,m}| + |a_{m-1,m-1}|) \end{aligned} \quad (7)$$

where we test whether or not the elements which appear in the positions $(m+1, m)$, $(m+2, m+1)$ are negligible compared with the three local diagonal elements $a_{m+1,m+1}$, $a_{m,m}$ and $a_{m-1,m-1}$.

Here we take m to be the largest integer ($\geq \ell$) for which condition (6) is satisfied.

For this computation, the mode bits are turned on for PE ℓ through PE n . The p_m , q_m and r_m for $m = \ell, \ell - 1, \dots, n$ are computed according to (6) on all PE's whose mode bits are turned on, and comparison is made to see whether (7) is satisfied. If (7) is satisfied on PE m , 1 is placed in the m th bit of the ACAR. Then the search is made for the lowest bit of the ACAR which contains 1, and m is set equal to this bit number. If no 1 is found in the ACAR, m is taken to be ℓ .

5.4 Double QR-transformation

A_{k+2} is computed by applying the QR-double transformation to A_k in such a way that $A_{k+2} = N_n^* \dots N_m^* A_k N_m \dots N_n$

where
$$N_i^* = I - \frac{U_i U_i^*}{2K_i^2} \quad \text{and} \quad U_i = (p_i \pm t_i, q_i, r_i, 0 \dots 0).$$

Here

$$\begin{aligned} p_i &= a_{ii}^2 - a_{ii} (s_k + s_{k+1}) + s_k s_{k+1} + a_{i,i+1} a_{i+1,i} \\ q_i &= a_{i+1,i} (a_{ii} + a_{i+1,i+1} - s_k - s_{k+1}) \\ r_i &= a_{i+2,i+1} a_{i+1,i} \quad \text{for } i = m \end{aligned}$$

and

$$\begin{aligned} p_i &= a_{i,i-1} \pm t_i, \quad q_i = a_{i+1,i-1} \\ \text{and } r_i &= a_{i+2,i-1} \quad \text{for } i \neq m. \end{aligned}$$

t_i and $2K_i^2$ are defined as

$$\begin{aligned} t_i &= \pm \sqrt{p_i^2 + q_i^2 + r_i^2} \\ 2K_i^2 &= t_i^2 + p_i t_i. \end{aligned}$$

Row modification:

For $i = m, m+1, \dots, n$, the elements of $N_i^* A_k^{(i)}$ are different from those of $A_k^{(i)}$ in only three rows, i.e., i th, $(i+1)$ th and $(i+2)$ th rows. These new elements are computed in the following way with the elements of $A_k^{(i)}$ denoted by a_{hj} :

$$\begin{aligned} (i,j) \text{ - element} &= a_{ij} - [(p_i \pm t_i) a_{ij} + q_i a_{i+1,j} + r_i a_{i+2,j}] \frac{1}{t_i} \\ (i+1,j) \text{ - element} &= a_{i+1,j} - [(p_i \pm t_i) a_{ij} + q_i a_{i+1,j} + r_i a_{i+2,j}] \frac{q_i}{2K_i^2} \\ (i+2,j) \text{ - element} &= a_{i+2,j} - [(p_i \pm t_i) a_{ij} + q_i a_{i+1,j} + r_i a_{i+2,j}] \frac{r_i}{2K_i^2} \end{aligned}$$

for $j = i, i+1, \dots, n$

In the actual computation, t_i and $2K_i^2$ are first computed and then p_i , q_i and r_i are found and stored in ADB's. The mode bits of PE's which contain $a_{i,i}$, $a_{i,i+1}$, \dots , $a_{i,n}$ are turned on and $c_i \equiv (p_i \pm t_i)a_{ij} + q_i a_{i+1,j} + r_i a_{i+2,j}$ are computed on these PE's. The computation of new elements

$$(i,j) - \text{element} = a_{ij} - \frac{c_i}{t_i}$$

$$(i,j+1) - \text{element} = a_{i,j+1} - \frac{c_i q_i}{2K_i^2}$$

$$(i,j+2) - \text{element} = a_{i,j+2} - \frac{c_i r_i}{2K_i^2}$$

are then performed.

Column Modification:

Similarly $(N_i^* A_k^{(i)})_{N_i}$ is computed from $N_i^* A_k^{(i)}$ for $i = m, m+1, \dots, n$ in the following way where the element of matrix $N_i^* A_i^{(k)}$ are denoted as $a_{j,h}$:

$$(j,i) - \text{element} = a_{ji} - [(p_i \pm t_i)a_{ji} + q_i a_{j,i+1} + r_i a_{j,i+2}] \frac{1}{t_i}$$

$$(j,i+1) - \text{element} = a_{j,i+1} - [(p_i \pm t_i)a_{ji} + q_i a_{j,i+1} + r_i a_{j,i+2}] \frac{q_i}{2K_i^2}$$

$$(j,i+2) - \text{element} = a_{j,i+2} - [(p_i \pm t_i)a_{ji} + q_i a_{j,i+1} + r_i a_{j,i+2}] \frac{r_i}{2K_i^2}$$

for $j = \ell, \dots, \min[i+3, n]$.

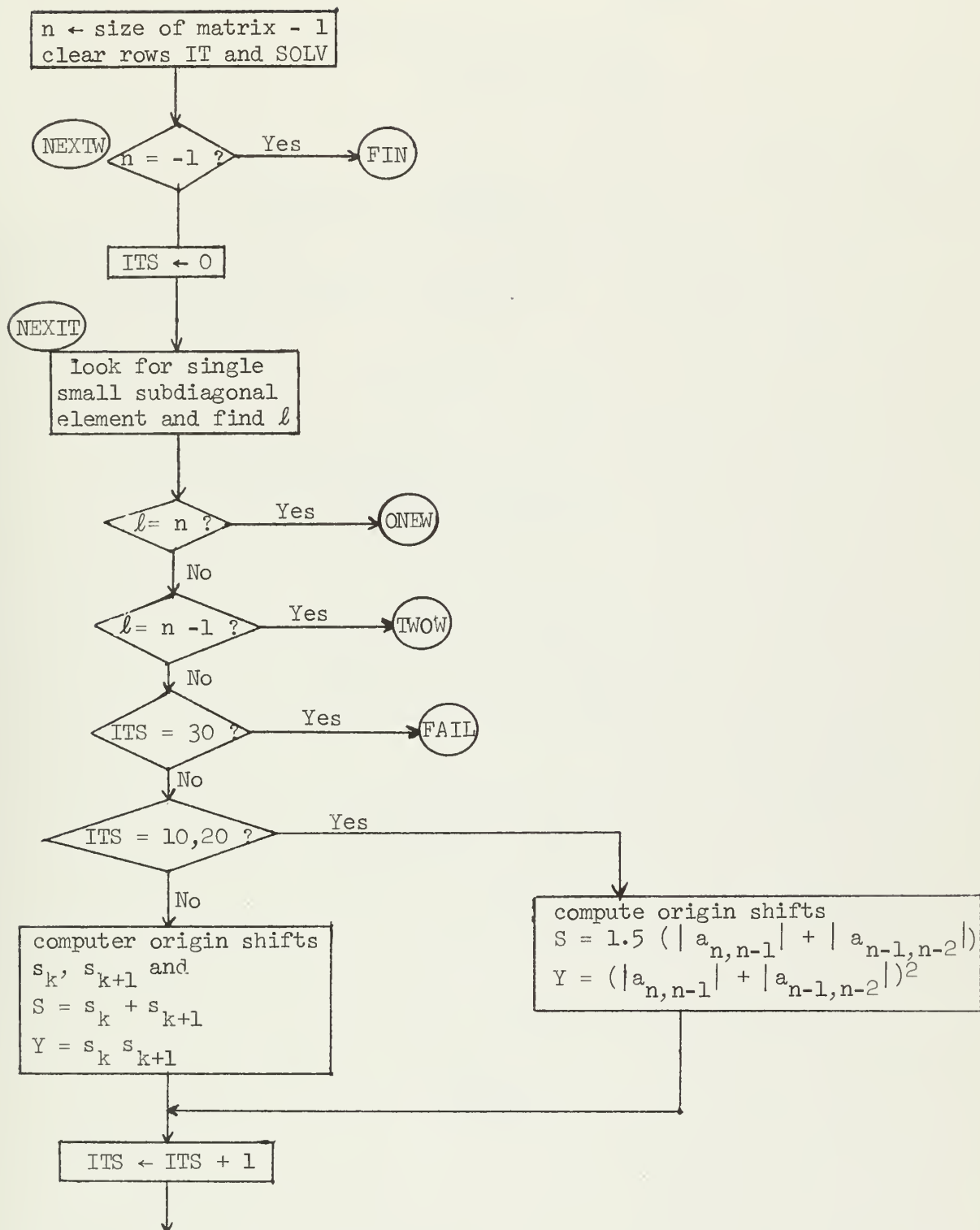
As in the row modification, $c_i' = (p_i \pm t_i)a_{ji} + q_i a_{j,i+1} + r_i a_{j,i+2}$ are computed on the PE's which contain a_{ji} for j from ℓ through $\min[i+3, n]$, then the computations of new $(j,i) - \text{element}$, $(j, i+1) - \text{element}$ and $(j, i+2) - \text{element}$ are performed on the corresponding PE's.

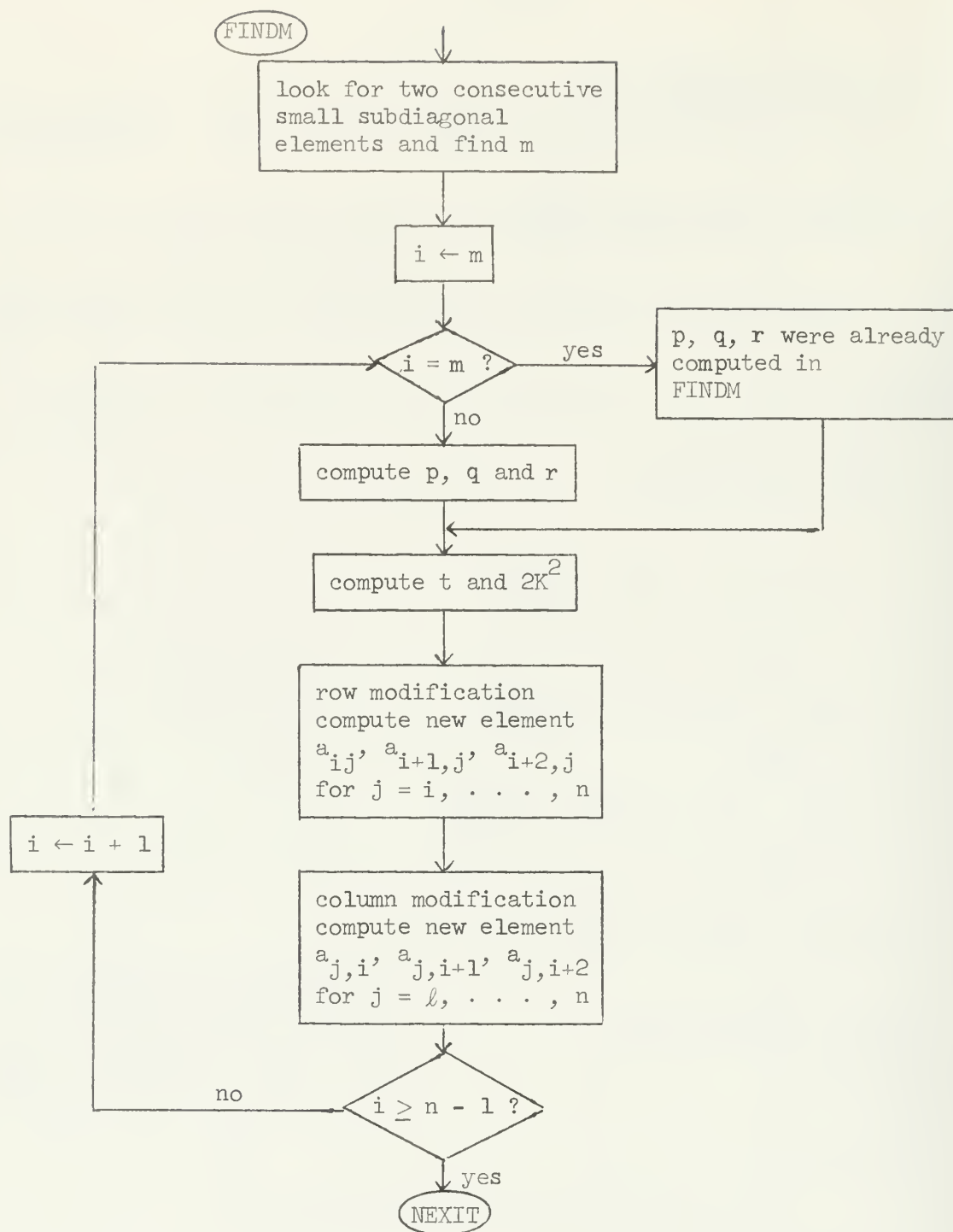
5.5 Computation of eigenvalues

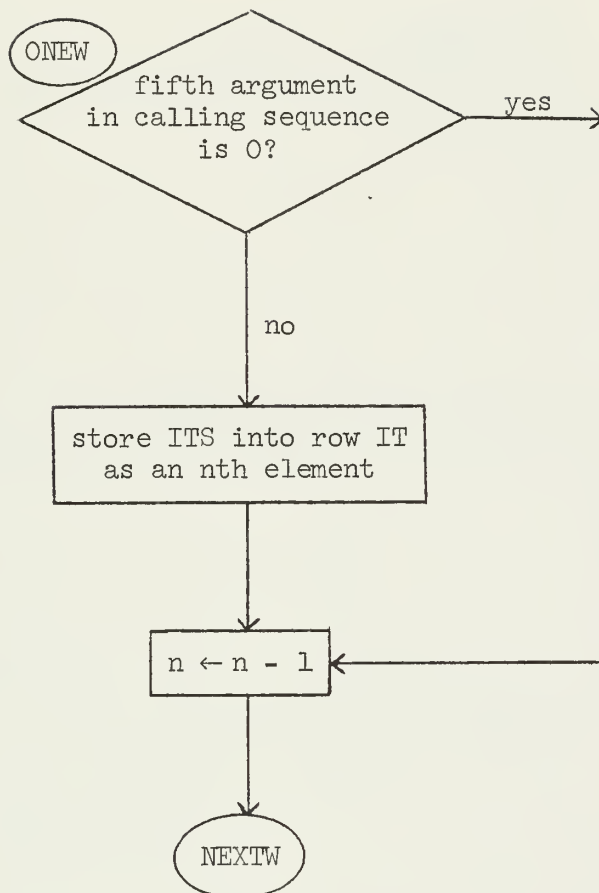
The eigenvalues are calculated as the last step of program after $a_{i+1,i}$ or $a_{i,i-1}$ $a_{i+1,i}$ become negligibly small for all $0 < i < n-2$. At each

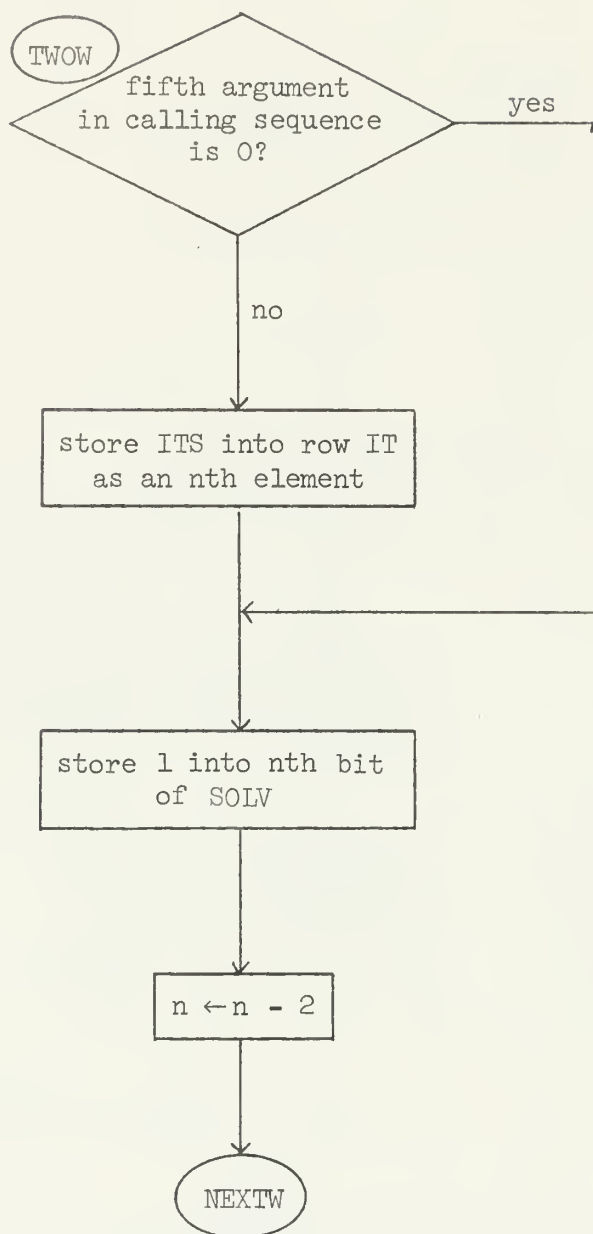
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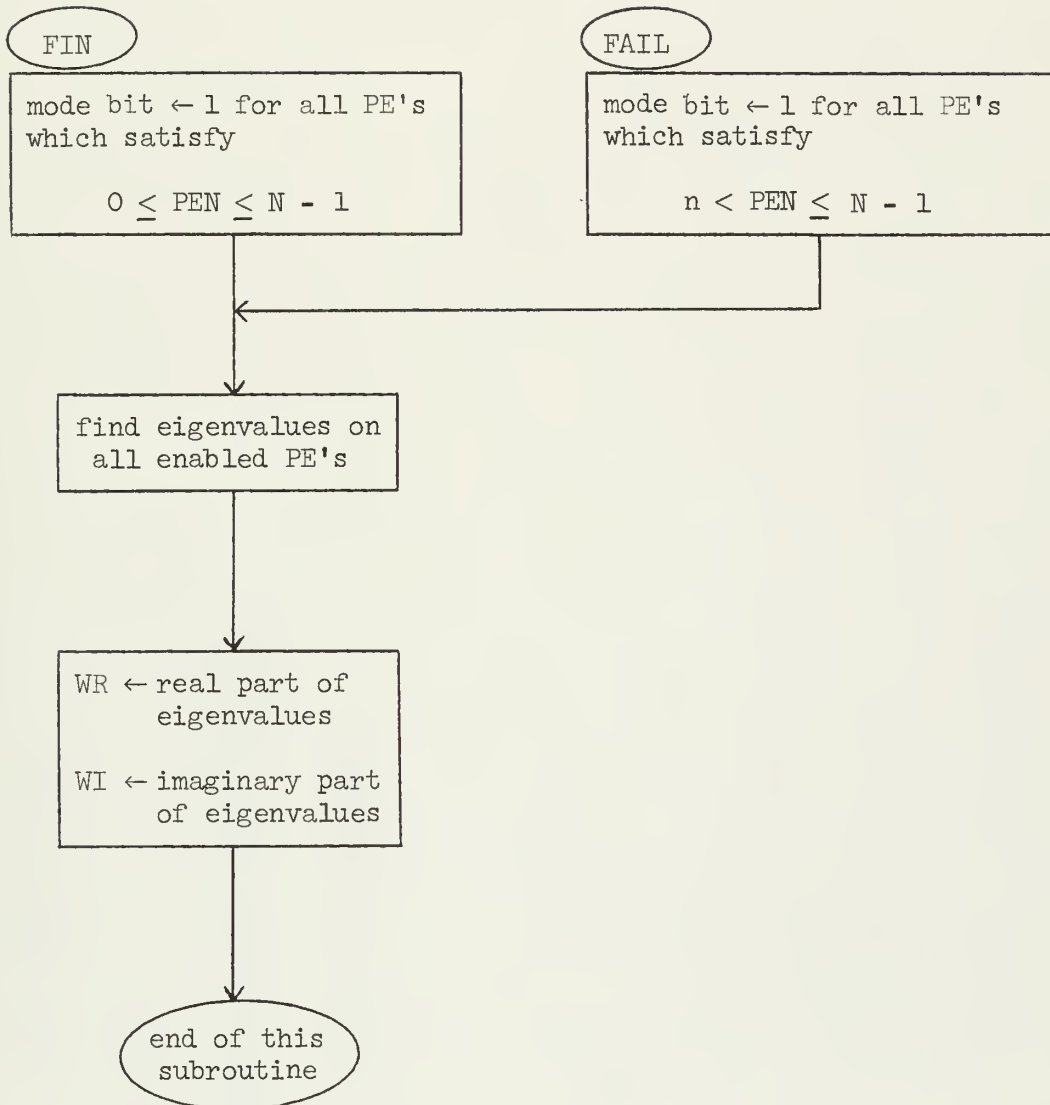
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3. B.N. Parlett: The LU and QR Algorithm. Mathematical Methods for Digital Computers, Vol. 2, A. Ralston and H. Wilf, Editors, J. Wiley, 1968.
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[illegible]


```

STL(1)      .ONE;
CSUB(0)     $C1;
STL(0)      .N;
ALIT(2)     =1;
LOAD(2)     $C0;
CSHR(0)     6;
STL(0)      .H;
ALIT(2)     =1;
LOAD(2)     $C0;
CSHR(0)     6;
STL(0)      .WR;
ALIT(2)     =1;
LOAD(2)     $C0;
CSHR(0)     6;
STL(0)      .WI;
ALIT(2)     =1;
LOAD(2)     $C0;
CSHR(0)     6;
STL(0)      .IT;
INITIALIZATION
LIT(3)      =00000000FFFFFFFF;
STL(3)      .MASK;
CLC(1)      ;
STL(1)      .SOLV;
LDL(0)      .N;
% BEGINNING OF N-LOOP, ACARO CONTAINS N
NEXTI: LIT(3) 0.0,-1;
CSUB(3)     $C0;
SETE        E OR -E;
SETE1       E OR E;
ZERF(3)     .NEXW;
JUMP        FIN;
NEXW: LIT(2)  =1.30.0;
STL(2)      .ITS;
% ITERATION COUNT .ITS IS SET TO 0
% DOUBLE OR ITERATION LOOP
% LOOK FOR SINGLE SMALL DIAGONAL ELEMENT
NEXTI: LDX
IXL
JXE
SETI        I AND E;
SETH        J OR E;
SETE        E OR -E;
SETE        H AND E;
SETE1       E OR E;
LDL(3)      .H;
LDA         +0(3);
SAP         J
RTL         $A.1;
ADRN        $R;
SLIT(3)     EPS;
LOAD(3)     $C3;
HLRN        $C3;
LDS         $A;
LDL(3)      .H;
% .ONE=1
% .N=(SIZE OF MATRIX)-1
%HI=PE ADDRESS OF H
% .WR=PE ADDRESS OF WR
% .WI=PE ADDRESS OF WI
% .IT=PE ADDRESS OF IT
% CLEAR CONTENT OF .SOLV
%ACARO=(SIZE OF MATRIX)-1
% ACAR3=-1-N
% IF N=-1 GO TO FIN
% ITERATION COUNT .ITS IS SET TO 0
% LOAD PE NUMBER TO RGX
%HI=1 IF RGX LEQ N
%EI=1 IF RGX LEQ N
$A:=H(L,L)
% TAKE ABSOLUTE VALUE
$R:=H(L-1,L-1)
% ACAR3=EPS
%RGSI=EPS+(ABS(H(L-1,L-1)))
00000056
00000057
00000058
00000059
00000060
00000061
00000062
00000063
00000064
00000065
00000066
00000067
00000068
00000069
00000070
00000071
00000072
00000073
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00000092
00000093
00000094
00000095
00000096
00000097
00000098
00000099
00000100
00000101
00000102
00000103
00000104
00000105
00000106
00000107
00000108
00000109

```

```

LDA      *1(3)I
SAP      I
RTL      $A,1I
LDA      $R,I
SBRN     $S,I
ISN      I
LIT(1)   =-1,1,0I
LUL(3)   .MASKI
CANO(3)  $C,I
CUR(1)   $C,I
SETC(2)  I
BT0      CTSBT(2)
TXGTM(1) .HT0I
CLC(1)   I

CONT1:   LDL(3)   .MASKI
          CANO(1)  $C,I
          CAND(3)  $C,I
          CSUB(3)  $C,I
          ZERT(3)  .ONEI
          CSUB(3)  .ONEI
          ZERT(3)  .TWOI
          SKIP
ONEW :   LUL(1)   .ITI
          ZERT(1)  .ONEI
          LUL(2)   .ITSI
          LUL(3)   .MASKI
          CANO(2)  $C,I
          SETE     E OR -EI
          LDX      PENI
          IXE      =0(0)I
          SETE     I AND EI
          SETE1    E OR EI
          LDA      $C2I
ONEW1:   LIT(3)   0,0,-1I
          CAD0(0)  $C3I
          JUMP     NEXTW
TWOW :   LUL(1)   .ITI
          ZERT(1)  .TWOI
          LUL(2)   .ITSI
          LUL(3)   .MASKI
          CANO(2)  $C3I
          SETE     E OR -EI
          LDX      PENI
          IXE      =0(0)I
          SETE     I AND EI
          SETE1    E OR EI
          LDA      $C2I
TWOW1:   LUL(2)   .SOLVI
          CSB(2)   0(0)I
          STL(2)   .SOLVI
          LIT(3)   0,0,-2I
          CAD0(0)  $C3I

          * LDA      *H+1
          * RGA:=ABS(H(L,L-1))

          * I:=1 IF RGA IS NEGATIVE

          * ACAR3:=0,0,N
          * ACAR1:=-1,1,N
          * ACAR2:=I
          * SKIP IF N-TH BIT OF $C2 IS 1
          * $C1:=0 IF NO SUBOIAE ELE IS SMALL
          * $C1:=L
          * $C3:=N
          * $C3:=N-L
          * IF L=N, GO TO ONEW
          * IF L=N-1, GO TO TWOW

          * ONE ROOT IS FOUND
          * IF IT=0, NO ITERATION COUNT
          * IS RECORDED
          * $C2:=ITERATION COUNT

          * IF IT IS NOT 0, IT:=ITERATION COUNT
          * DECREASE ACARO BY 1
          * AND GO TO NEXTW
          * TWO ROOTS ARE FOUND
          * IF IT=0, NO ITERATION COUNT
          * IS RECORDED
          * $C2:=ITERATION COUNT

          * IF IT IS NOT 0, IT:=ITERATION COUNT
          * N-TH HIT OF .SOLV:=1
          * DECREASE ACARO BY 2

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00000165      JUMP NEXTW;
AFTJ : CAND(1) .MASKI;
        STL(1) .L;
        LDL(2) .ITS;
        TXLF(2) SC2=FAILJ;
CAND(2) .MASK1;
        LIT(1) =10;
        CSUB(2) SC1;
        ZERT(2) .IT10;
        CSUB(2) SC1;
        ZERT(2) .IT10;
        ZERT(2) .ITN10;
        SKIP FAIL;
X FORM SHIFT S AND Y WHEN IT=10,20
IT10 : LDL(3) .H;
        CADD(3) SC0;
        LDA O(3);
        SAP ;
        LDS $A;
        LDA -1(3);
        SAP ;
        RTL $A,1;
        LDA $R;
        ADRN $S;
        LDS $A;

X          XRGSI=ABS(H(N,N-1))
          +ARS(H(N-1,N-2))

        LIT(3) =1,5;
        MLRN SC3;
        IXE --1(0);
        SETE I AND E;
        SETE E OR E;
        LDC(2) $A;
        STL(2) .S;
        LDA $S;
        MLRN $S;
        LDC(3) $A;
        STL(3) .Y;
        SKIP ,CONT15;
X FORM SHIFT S AND Y WHEN IT NEQ 10,20
ITN10 : LDL(3) .H;
        CADD(3) SC0;
        LDA O(3);
        LDS -1(3);
        RTL $S,1;
        IXE =0(0);
        SETG I AND E;
        SETE I AND E;
        SETE1 E OR E;
        ADRN $R;
        LDC(2) $A;
        STL(2) .S;
        SETE E OR -E;
        SETE H AND E;
        SETE1 E OR E;

X          X.S:=H(N,N)+H(N-1,N-1)
          XE1=1 IF RGX LEQ N

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00000219      LDLC(3)      .H)
00000220      CADD(3)      $C0)
00000221      LDA          0(3))
00000222      MLRN
00000223      LDS          $A)
00000224      LDA          0(3))
00000225      RTL          $A,1)
00000226      LDA          $R)
00000227      SETE        G AND E)
00000228      SETE1       E OR E)
00000229      MLRN        -1(3))
00000230      CHSA
00000231      ADRN
00000232      LDC(3)     $A)
00000233      STLC(3)     .Y)
00000234
00000235      * LOOK FOR TWO CONSECUTIVE SMALL DIAGONAL ELEMENT FOR M=N-2 TO L
00000236      CONT5(LDL(1) .ITS)
00000237      LDLC(2)      .L)
00000238      ALIT(1)      =1)
00000239      STLC(1)      .ITS)
00000240      JXG          =0(2))
00000241      IXL          =-1(0))
00000242      SETE         E DR -E)
00000243      SETE         J AND E)
00000244      JXE          =0(2))
00000245      SETE         J DR E)
00000246      SETG         I AND E)
00000247      * COMPUTE P=H(M,M)+(H(M,M)-S)+Y+H(M+1,M)+H(M,M+1)
00000248      SETE         E DR -E)
00000249      SETE         H AND E)
00000250      SETE1        E OR E)
00000251      LDLC(3)      .H)
00000252      LDA          +0(3))
00000253      LDS          $A)
00000254      LDLC(3)      .S)
00000255      SBRN         $C3)
00000256      SETE         G AND E)
00000257      SETE1        E OR E)
00000258      MLRN
00000259      LDS          $S)
00000260      SETE         E DR -E)
00000261      SETE         H AND E)
00000262      SETE1        E OR E)
00000263      LDLC(3)      .H)
00000264      LDA          +-1(3))
00000265      RTL          $A,-1)
00000266      LDA          $R)
00000267      SETE         G AND E)
00000268      SETE1        E DR E)
00000269      MLRN        +1(3))
00000270      ADRN         $S)
00000271      LDLC(2)      .Y)
00000272      ADRN         $C2)
00000273      STA          PI
00000219      * LDA      H(0)
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SETI  E OR -E;
SETI  H AND E;
SETI  E OR E;
LDA  Q;
SAP  ;
LDS  $A;
LDA  R;
SAP  ;
ADRN  $S;
LDS  $A;
H;
LDA  +1(3);
SAP  ;
RTL  $A,1;
LDA  L;
LDL(2)  =0(2);
JXE  G AND E;
SETI  -J AND E;
SETI  E OR E;
SETI  E OR E;
LDA  $R;
MLRN  $S;
LDL(3)  =R;
STA  0(3);
SETI  E OR -E;
SETI  H AND E;
SETI  E OR E;
H;
LDL(3)  +0(3);
SAP  ;
LDS  $A;
RTL  $S,1;
SETI  I AND E;
SETI  E OR E;
ADRN  $R;
RTL  $S,-1;
ADRN  $R;
LDS  $A;
LDA  P;
SAP  ;
SLIT(2)  EPS;
LOAD(2)  $C2;
MLRN  $C2;
MLRN  $S;
CHSA  ;
LDL(3)  =R;
ADRN  0(3);
JSN  ;
LIT(1)  0,0,-2;
CADD(1)  $C0;
CAND(1)  =MASK1;
LIT(3)  =-1,0,0;
COR(1)  $C3;
LDL(3)  =L;
CADD(3)  =ONE;

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CSHL(3) 241
COR(1) 303
SETC(3) J
BT1 : CTSBT(3) 0(1),CONT2;
TXGTM(1) ,RT1;
LDL(1) ,LJ
CONT2: CAND(1) ,MASK1;
STL(1) ,HJ
X PUT ZERO IN APPROPRIATE PLACES
SETE E OR -E;
SETE1 E OR E;
CLRA
IXG =0(1);
JXE =0(1);
SETE I AND E;
SETE J OR E;
JXL =-2(0);
SETE J AND E;
SETE1 E OR E;
LDL(3) ,HJ
STA +3(3);
IXE =-2(0);
SETE I OR E;
SETE1 E OR E;
STA +2(3);
X PUT P,Q,R INTO ADB WHEN K=M
SETE E OR -E;
SETE1 E OR E;
CLRA
IXE =0(1);
SETE I AND E;
SETE1 E OR E;
LDA P;
LDS Q;
LDC(2) ,SA;
LDC(3) ,SS;
STL(2) ,P;
STL(3) ,Q;
LDA R;
LDC(2) ,SA;
STL(2) ,R;
X ARRANGE MATRIX INTO SKEW STORAGE FORM
LIT(2) -1,1,0;
COR(2) ,SC0;
SETE E OR -E;
SETE1 E OR E;
DTAR1: LDL(3) ,HJ
CAOD(3) ,SC2;
LDA 0(3);
RTL SA,0(2);
STR 0(3);
TXGTM(2) ,DTAR1;
X DOUBLE QR STEP INVOLVING ROWS L TO N AND COLUMNS M TO N
LIT(3) =0,1,-1;
X ACAR3:=0,L+1,0
X ACAR1:=-1,L+1,N-2
XACAR3:=1 IF J CONTAINS 1
X ACAR1:=0,0,M
X,M:=0,0,M
X E,E1:=1 FOR M LEQ RGX < N-2
X STA +H+3
X E,E1:=1 FOR M LEQ RGX LEQ N-2
X STA +H+2
SPICK UP M-TH ELEMENT OF
X P,Q AND R WHEN K=M
X I:=1 IF RGX=M
X E,E1:= 1
X,P:=M-TH ELEMENT OF P
X,Q:=M-TH ELEMENT OF Q
XACAR2:=1,1,1,N
X LDA M(2)
X STR M(2)
X INITIALIZATION OF K-LOOP

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CADD(3)  $C01
CSHL(3)  241
COR(1)   $C31
% ACAR1 CONTROLS K=LOOP, KIM STEP 1 UNTIL N=1
NEXTK: LCL(2)  $C01
CSUB(2)   0NE1
CSUB(2)   $C11
ZERT(2)   1LAST1
LCL(3)    0NE1
STL(3)    1NOTLAST1
SKIP
LAST: CLC(3)
STL(3)    1NOTLAST1
CONT17:
    SETE   F OR -E1
    SETE1  E OR E1
    CLRA   1
    LCL(3) 1MASK11
    CAND(3) $C11
    CSUB(3) 1M1
    ZERT(3) 1KEQM1
% COMPUTE P=H(K,K-1),Q=H(K+1,K-1) AND R=H(K+2,K-1) WHEN K NEQ N
LDX
RTL
$X,-1(1)1
LDX
$R1
IXE
=2(1)1
SETH
I AND E1
JXE
=1(1)1
IXE
=0(1)1
SETE
I AND E1
SETE
J OR E1
SETG
H OR E1
SETE
E OR -E1
SETE
G AND E1
SETE1
E OR E1
LCL(3)
1M1
LOA
+0(3)1
1NOTLAST1
LDL(2)
1CPQR1
ZERF(2)
SETE
H AND E1
SETE1
E OR E1
CLRA
1
SETE
E OR -E1
SETE
G AND E1
SETE1
E OR E1
$A1
LOS
SAP
RTL
AORN
$R1
$A,21
RTL
ADRN
$R1
SETE
H AND E1
SETE1
E OR E1
$ACAR3:=0,1,N-1
$ACAR3:=1,N-1,0
$ACAR1:=1,N-1,M
% SET NOTLAST WITH APPROPRIATE
% BOOLEAN EXPRESSION
$ACAR2:=(N-1)-K
% NOTLAST:=1 IF K NEQ N-1
%NOTLAST:=0 IF K=N-1
% TEST IF K=M
$ACAR3:=0,0,K
% ,ACAR3:=K-M
% IF K=M GO TO KEQM,
% P,Q,R ARE ALREADY COMPUTED
% HI=1 FOR H(K+2,K-1)
% JI=1 FOR H(K+1,K-1)
% II=1 FOR H(K,K-1)
% GI=1 FOR H(K,K-1),H(K+1,K-1)
% AND H(K+2,K-1)
% E,EI:=G
% LDA +H
% IF NOTLAST=0, H(K+2,K-1)
% IS CONSIDERED TO BE 0
% RESET E,EI
$RGS:=H(K,K-1),H(K+1,K-1)
% AND H(K+2,K-1)
% RGA1=ABS(RGA)
$RGA1=ABS(H(K,K-1))+ABS(H(K+1,K-1))
% +ABS(H(K+2,K-1)) AT M=1
$EI=M

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LDC(2)      SAJ
STL(2)      .XXJ
ZERF(2)     .NXP0J
JUMP        CONT3J
NXPOJ:      E OR -EJ
            SETE  G AND EJ
            SETE1 E OR EJ
            CLRA  J
            LDB  SAJ
            LDA  SSJ
            DVM  SC2J
            LDS  SAJ
            .CONT27J
            SKIP
            $ PLACE P,Q,R IN APPROPRIATE REGISTER WHEN K=M
KEQM:      LDX  PENJ
            LDL(2) .MASK1J
CAND(2)    SC1J
ZERT(2)    .KZEROJ
RTL        SX,-1(1)J
SKIP
KZEROJ:    RTL  .KEQM2J
KEQM2J:    LDX  SX,0(1)J
            SRJ
            .2(1)J
            I AND EJ
            JXJ
            IXJ
            SETE  J AND EJ
            SETE  H OR EJ
            SETE  E OR -EJ
            .PJ
            .OJ
            I AND EJ
            SETE1 E OR EJ
            ADRN  SC2J
            SETE  E OR -EJ
            SETE  J AND EJ
            SETE1 E OR EJ
            ADRN  SC3J
            .PJ
            .OJ
            E OR -EJ
            SETE  H AND EJ
            SETE1 E OR EJ
            ADRN  SC2J
            SETE  E OR -EJ
            SETE1 E OR EJ
            LDS  SAJ
            $ COMPUTE S=SQRT(P**2+Q**2+R**2)
CONT27J:   SETC(2)  JI
            STL(2)  .JBITI
            JSN     J
            SETE    I AND EJ
            SETJ    J AND EJ

SACAR2:=RGA AT M=1
$ .XXJ:=ACAR2
$ IF .XX=0 GO TO CONT3
$ E,E1:=G

SRGS:=P/XX.Q/XX.R/XX

$ DUMMY BIT SETTING FOR COMPUTING
$ S IN CONT27
SI:=1 FOR H(K+2,K)
SJ:=1 FOR H(K+1,K)
SI:=1 FOR H(K,K)

$ GI:=1 FOR H(K,K),H(K+1,K),H(K+2,K)

$ E,E1:=I
SRGA:=0,...0,P,0,0,...0
$ E,E1:=J
SRGA:=0,...0,P,0,0,...0
$ E,E1:=H
SRGS,RGA:=0,...0,P,Q,R,0,...0

ISAVE J BIT INTO .JBIT
SI:=SIGN BIT OF RGA
SI:=1 IF P IS NEGATIVE

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SETI      J AND E1
SETI      E OR E1
LDC(2)    $A1
STL(2)    .Y1
SETI      E OR -E1
SETI      H AND E1
SETI      E OR E1
LDC(3)    $A1
STL(3)    .Z1
* COMPUTE Q=Q/P AND R=R/P FOR ROW AND COLUMN MODIFICATION
SETI      E OR -E1
SETI      J AND E1
SETI      H OR E1
SETI      E OR E1
LDC(3)    .P1
CLRA      J
LDB        $A1
LDA        $S1
DVRN      $C31
SETI      J AND E1
SETI      E OR E1
LDC(2)    $A1
STL(2)    .Q1
SETI      E OR -E1
SETI      H AND E1
SETI      E OR E1
LDC(3)    $A1
STL(3)    .R1
* COMPUTE NEW H[K,K-1]
* IF K NEO M THEN H[K,K-1] := -S+X
* ELSE IF L NEO M THEN H[K,K-1] := -H[K,K-1]
LDC(2)    .MASK1
CAND(2)   $C11
CSUB(2)   .M1
ZERT(2)   .CPLM1
SETI      E OR -E1
SETI      I AND E1
SETI      E OR E1
CLRA      J
LDC(2)    .S1
LDC(3)    .XX1
SBRN      $C21
MLRN      $C31
LDC(3)    .H1
STA        *O(3)1
SKIP      .RWMOD1
CPLM      LDC(2)
          .L1
          CSUB(2)
          ZERT(2)
          LDX
          RTL
          LDX
          IXE
          SETI
          E OR -E1
          PEN1
          $X,-1(1)1
          $R1
          =O(1)1
          E OR -E1
          $H(K,K-1) := -H(K,K-1) IF L NEO M, K=M
          $ ACAR1 CONTAINS M WHICH IS
          GREATER THAN L
          $ STA *H
          $ TEST IF L=M
          $H(K,K-1) := -H(K,K-1) IF L NEO M, K=M
          $ ACAR1 CONTAINS M WHICH IS
          GREATER THAN L
          $ IF K NEO M, H[K,K-1] := -S+X
          $ E, E1 := I
          $ BETWEEN L AND M
          $ IF K=M SKIP TO COMPARISON
          $ ACAR2 := 0, 0, K
          $ ACAR2 := 0, 0, K-M
          $ ACAR2 := 0, 0, K
          $ R1 := R/.P
          $ Q1 := Q/.P
          $ RGA1 := Q/.P, .P, .R/.P FOR J=1 AND M=1
          $ E, E1 := J OR M
          $ Z1 := R/.S
          $ Y1 := Q/.S
          $ E, E1 := H
          $ E, E1 := J

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```

SETI I AND E1
SETI1 E OR E1
LDL(3) *H1
LDA +0(3)1
CHSA 1
STA +0(3)1

*
DEFINE KRGXN =
  IXG =0(1)1
  JXL =0(0)1
  SETI I AND E1
  SETI J AND E1
  IXE =0(1)1
  JXE =0(0)1
  SETI I OR E1
  SETI J OR E11

* RMOD1 SETI E OR -E1
* ROW MODIFICATION
LDX PEN1
RTL $X,2(1)1
LDX $R1
KRGXN 1
SETH E OR E1
SETI E OR -E1
RTL $X,-1)1
LDX $R1
KRGXN 1
SETJ E OR E1
SETI E OR -E1
SETI1 E OR E1
LDL(3) *H1
CADD(3) $C11
LDA 1(3)1
SETI G AND E1
SETI1 E OR E1
LDL(2) *J1
MLRN $C21
LOS $A1
SETI E OR -E1
SETI1 E OR E1
LDA 0(3)1
RTL $A,11
LDA $R1
ADRN $S1
LOS $A1
LDL(2) *NOTLAST1
ZERT(2) *CONTA41
* IF NOTLAST=1 COMPUTE
* H[K+2,J]1= H[K,J]+0+H[K+1,J]+R+H[K+2,J])1*Z

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LDA 2(3);
RTL $S,1;
LDS $R;
SETL H AND E;
SETL E OR E;
LDL(2) .R;
LDL(3) .Z;
MLRN $C2;
ADRN $S;
LDS $A;
MLRN $C3;
CHSA J;
LDL(3) .H;
CADD(3) $C1;
ADRN 2(3);
STA 2(3);
SETL E OR -E;
SETL E OR E;
RTL $S,-1;
LDS $R;

X COMPUTE
X H(K+1,J):= H(K,J)+Q*H(K+1,J)+R*H(K+2,J)*Y
CONT4: SETL E OR -E;
SETL G AND E;
SETL E OR E;
LDL(2) .Y;
LDA $S;
CHSA J;
MLRN $C2;
LDL(3) .H;
CADD(3) $C1;
ADRN 1(3);
STA 1(3);

X COMPUTE
X H(K,J):= H(K,J)-(H(K,J)+Q*H(K+1,J)+R*H(K+2,J))*X
SETL E OR -E;
SETL J AND E;
SETL E OR E;
RTL $S,-1;
LDL(3) .X;
LDA $R;
CHSA J;
MLRN $C3;
LDL(3) .H;
CADD(3) $C1;
ADRN 0(3);
STA 0(3);

X COLUMN MODIFICATION
LDL(2) ,MASK1;
LIT(3) 0,0,3;
CAND(2) $C1;
CADD(2) $C3;
LDL(3) $C2;

XRGAI=H+2(1)
X ROUTE RGS BY 1

X E*E1:=H

X RGS:=RGA+.R*RGs
X RGA:=.Z*RGs

X ADRN H+2(1)
X H+2(1):=H+2(1)+RGA

X ROUTE RGS BY -1

X E*E1:=G

X RGA:=-RGS+.Y

X ADRN H+1(1)
X STA H+1(1)

X E,E1:=J
X ROUTE RGS BY -1

XRGAI=-RGS+.X

X ADRN H(1)
X H(1):=H(1)+RGA

X SET J WITH K+3 OR N
X ACAR2:=0,0,K
X ACAR2:=0,0,K+3
X ACAR3:=0,0,K+3

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00000712
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00000763

```

CSUB(2)      $C01
CTSRF(2)     40,NINTUJ1
STL(3)       .J1
SKIP         .CONTS1
NINTUJ1:STL(0) .J1
*
DEFINE LRGXJ =
LDL(2)       .L1
LDL(3)       .J1
IXG         =0(2)1
JXL         =0(3)1
SETE        E OR =E1
SETE        I AND E1
SETE        J AND E1
IXF         =0(2)1
JXF         =0(3)1
SETE        I OR E1
SETE        J OR E 1*1
*
CONTS1: SETE E OR =F1
LDX         PFN1
RTL         $X,0(1)1
LDX         $R1
LRGXJ       I
SETH        E OR E1
SETE1       E OR E1
LDL(2)      .X1
LDL(3)      .-1
LDA         +0(3)1
MLRN        $C21
LDS         $A1
SETE        E OR =F1
SETE1       E OR E1
RTL         $X,11
LDX         $R1
RTL         $S,11
LDS         $Q1
LRGXJ       I
LDL(2)      .Y1
SETG        E OR E1
SETE1       E OR E1
LDL(3)      .H1
LDA         +0(3)1
MLRN        $C21
ADRN        $S1
LDS         $A1
* COMPUTE H[I,K+2]=H[I,K+2]-(X*H[I,K]+Y*H[I,K+1]+Z*H[I,K+2])*R
*
SETE        E OR =E1
SETE1       E OR E1
LDL(2)      .NOTLAST1
ZERT(2)     .CONT61
RTL         $X,11
LDX         $R1
RTL         $S,11
*
* ACAR21=0,0,K+3-N
* IF 40-TH HIT OF ACAR2 IS 0,J1=N
* J1=K+3 IF K+3<N
* J1=N IF K+3 GEQ N
*
* I1=1 IF L<RGX
* J1=1 IF RGX<J
*
* E1=1 IF L<RGX<J
* I1=1 IF L=RGX
* J1=1 IF RGX=J
*
* E1=1 IF L LEQ RGX LEQ J
*
* RGX1=PEN
* ROUTE RGX BY K
*
* H IS SET FOR H[I,K]
*
* RGA1=H[I,K]
* RGS1=H[I,K]*X
*
* ROUTE RGX BY 1
* ROUTE RGS BY 1
*
* G IS SET FOR H[I,K+1]
*
* RGA1=H[I,K+1]
* $A1=H[I,K+1]*Y
*
* $S1=X*H[I,K]+H[I,K+1]
* $A1=Y*H[I,K+1]+Z*H[I,K+2])*R
* IF NOTLAST=1
*
* IF NOTLAST=0, GO TO CONT6
* NO NEED TO COMPUTE H[I,K+2]

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00000A73
00000A74
00000A75
00000A76
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CADD(3)  SC11
LDA      0(3)1
LIT(2)   =0.0.641
CSUB(2)  SC11
RTL      SA.0(2)1
STR      0(3)1
TXGTM(1) 0TAR2;
JUMP     NEXTIT1
% IF ALGORITHM FAILS.G1=1 WHERE SURDIAGONAL ELEMENTS CONVERGED
FAIL 1:SETE E OR =F1
LDX      PFN1
TXG      =0(0)1
LDL(3)   11
JXL      =0(3)1
SETE     I AND F1
SETE     J AND E1
JXE      =0(3)1
SETG     J OR E1
JUMP     ETG1
% IF PROCESS IS COMPLETED.G1=1 FOR ALL RITS
FIN 1:SETE E OR =F1
LDX      PFN1
TXG      =01
LDL(3)   11
JXL      =0(3)1
SETE     I AND F1
SETE     J AND F1
IXE      =01
JXE      =0(3)1
SETE     I OR E1
SETG     J OR E1
% FIND ONE ROOT
EIG 1:SETE E OR =F1
LDL(3)   11
IXL      =0(3)1
JXE      =0(3)1
SETE     I AND E1
SETE     J OR E1
SETE     G AND E1
SETE1    E OR E1
LDL(3)   11
LDA      =0(3)1
LDL(3)   11
STA      0(3)1
CLRA     1
LDL(3)   11
STA      0(3)1
% TWO ROOTS ARE FOUND
LDL(2)   11
ZERF(2)  11
JUMP     MOREND1
TWRT 1: SETE E OR =E1
TXG      =01
LDL(3)   11
JXL      =0(3)1
% LDA  =N
% JXL  =N
% STR  H(1)
% TEST AND MODIFY ACAR2 BY -1
% JXL  =N
% JXE  =N
% ACAR0 CONTAINS PRESENT MATRIX
% SIZE
% JXL  =N
% JXE  =N
%G1=1 FOR 0 LEQ RGX LEQ N
% IXL  =N
% JXE  =N
% LDA  =H
% STA  =R
% IMAGINARY PART= 0
% SET APPROPRIATE BIT PATTERN
% JXL  =N

```



```

LDL(2)      $C1;
CAND(1)     $C0;
CAND(1)     .SOLV;
ZERT(1)     .FREAL;
LDFE1      $C1;
LDL(3)      .W1;
STA         0(3);
LDS        TMP2;
LDL(3)      .WR;
STS        0(3);
CSHL(1)     1;
LDEE1      $C1;
RTL        $A*-1;
LDA        $H;
RTL        $S*-1;
CHSA       J;
LDL(3)      .W1;
STA         0(3);
LDL(3)      .R;
STR         0(3);

% FIND REAL PAIR
FREAL: COMPC(0) J
CAND(2)     $C0;
CAND(2)     .SOLV;
ZERT(2)     .HQREND;
LDEE1      $C2;
LDA        TMP2;
LDS        $A;
ADRN       TMP3;
LDL(3)      .R;
STA         0(3);
CLRA       J;
LDL(3)      .W1;
STA         0(3);
LDA        TMP3;
CHSA       J;
ADRN       $S;
CSHL(2)     1;
LDEE1      $C2;
RTL        $A*-1;
LDL(3)      .R;
STR         0(3);
CLRA       J;
LDL(3)      .W1;
STA         0(3);

HQREND:
LDL(0)      .HQSV3;
LDEE1      $C0;
LDL(0)      .HQSV0;
LDL(1)      .HQSV1;
LDL(3)      .HQSV3;
EXCHL(3)    $ICR;
END

% ACAR1:=H RIT AND SIGN
% ACAR1:=ACAR1 AND SLTN PATTERN

% STA W1
% FOUND WR(N) AND WI(N)

% STA WR
% SHIFT ACAR1 LEFT BY 1
% E,E1:=ACAR1
% ROUTE RGA BY -1
% ROUTE RGS BY -1

% STA W1
% STR W

% ACAR0:=COMPLEMENT OF ACAR0
% ACAR2:=H AND C(SIGN)

% STA WR
% STA W1

% RGA:=TMP2-TMP3

% STR WR

% STA W1

% RESTORE E,E1

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13. ABSTRACT The implementation of QR-algorithm on ILLIAC IV is described. An ASK subroutine for computing all eigenvalues of a real Hessenberg matrix of order less than or equal to 64 by this algorithm is attached. The QR-transformation consists of the decomposition of the matrix A_k into the product of a unitary matrix Q_k and an upper triangular matrix R_k , and forming A_{k+1} by post-multiplying R_k by Q_k , where $A_1 = A$ is the original matrix. All eigenvalues are either isolated on the diagonal or are eigenvalues of a 2 x 2 diagonal submatrix as $k \rightarrow \infty$.			

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